Logic I: Fast Lecture 05

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Readings refer to sections of the course textbook, *Language, Proof and Logic*.

1. Vegetarians Are Evil

Reading: §9.2, §9.3, §9.5

 $\forall x (Evil(x) \rightarrow HatesMeat(x))$

 $\forall x (HatesMeat(x) \rightarrow Vegetarian(x))$

 $\forall x (Vegetarian(x) \rightarrow Evil(x))$

2. Counterexamples with Quantifiers

	Evil(x)	HatesMeat(x)	Vegetarian(x)
Ayesha	no	no	yes

3. Something Is Above Something

Reading: §11.1 Something is above something: ∃x ∃y Above(x,y)

4. Multiple Quantifiers: Everyone Likes Puffins

Reading: §11.1 I like puffins: $\forall x (Puffin(x) \rightarrow Likes(a,x))$ y likes puffins: $\forall x (Puffin(x) \rightarrow Likes(y,x))$ Everyone likes puffins: $\forall y \forall x (Puffin(x) \rightarrow Likes(y,x))$

5. Relations: Transitivity

Reading: §15.1

A *transitive* relation is one such that if x bears it to y and y bears it to z then x bears it to z. (E.g. LeftOf is transitive; NotAdjacent is not transitive.)

6. Expressing Relations with Quantifiers

Reading: §15.1

A *reflexive* relation is one that everything bears to itself. (E.g. SameShape) reflexive: $\forall x R(x,x)$



A *symmetric* relation is one such that if x bears it to y, then y bears it to x. (E.g. Adjacent(x,y)) symmetric: $\forall x \forall y (R(x,y) \rightarrow R(y,x))$



A *transitive* relation is one such that if x bears it to y and y bears it to z then x bears it to z. (E.g. LeftOf is transitive; DifferentShape is not transitive)

transitive: $\forall x \forall y \forall z \ (\ (\ R(x,y) \land R(y,z) \) \longrightarrow R(x,z) \)$



7. Expressing Counterexamples Formally

Reading: §15.1

Give a counterexample to this argument:

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 \begin{array}{l} \forall x \ \mathsf{R}(x,x) \\ \hline \forall x \forall y \forall z \ ( \ [ \ \mathsf{R}(x,y) \land \mathsf{R}(y,z) \ ] \to \mathsf{R}(x,z) \ ) \\ \hline \forall x \forall y \ [ \ \mathsf{R}(x,y) \to \mathsf{R}(y,x) \ ] \end{array}
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Informally:



Formally:

Domain: {a, b} R: {<a,a>, <a,b>, <b,b>}

8. There Is a Store for Everything

Reading: §11.2, §11.3

There is a store for everything:

∃y∀x StoreFor(y,x)

∀y∃x StoreFor(x,y)

Other sentences to translate:

Wikipedia has an article about everything Everyone hurts someone they love Someone hurts everyone she loves

9. ∀Intro: An Incorrect Proof

Reading: §13.1, §13.2 This proof is wrong, but why?:



There is a counterexample to the argument:



10. Two Things Are Broken

Reading: §14.1

To translate sentences involving number into FOL, use identity. For example,

'Two things are broken' might be translated as:

 $\exists x \; \exists y \; (\; Broken(x) \land Broken(y) \land \neg(x{=}y) \;)$

11. Quantifier Equivalences: $\forall x(Square(x) \rightarrow Broken(x)) =$ $\forall x(\neg Broken(x) \rightarrow \neg Square(x))$

Reading: §10.3

ΡQ	$P \rightarrow Q \neg Q \rightarrow \neg P$
ТТ	т т
ΤF	F F
FΤ	тт
FF	тт

12. Quantifier Equivalences:
∀x Created(x) = ¬∃x Created(x)

Reading: §10.3, §10.4

13. Soundness and Completeness: Statement of the Theorems

Reading: §8.3, §13.4

'A \vdash B' means there is a proof of B using premises A

' \vdash B' means there is a proof of B using no premises

'A \models B' means B is a logical consequence of A

 $\doteq B'$ means B is a tautology

'A \models_{TT} B' means B is a logical consequence of A just in virtue of the meanings of truth-functions (the textbook LPL calls this 'tautological consequence')

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Soundness: If A \vdash B then A \models B
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i.e. if you can prove it in Fitch, it's valid

Completeness: If $A \models_{TT} B$ then $A \vdash B$

i.e. if it's valid just in virtue of the meanings of the truth-functional connectives, then you can prove it in Fitch.